

# Box-Ball Systems and Robinson–Schensted–Knuth Tableaux

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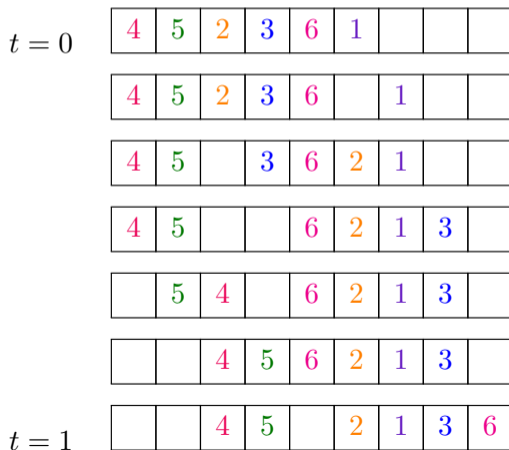
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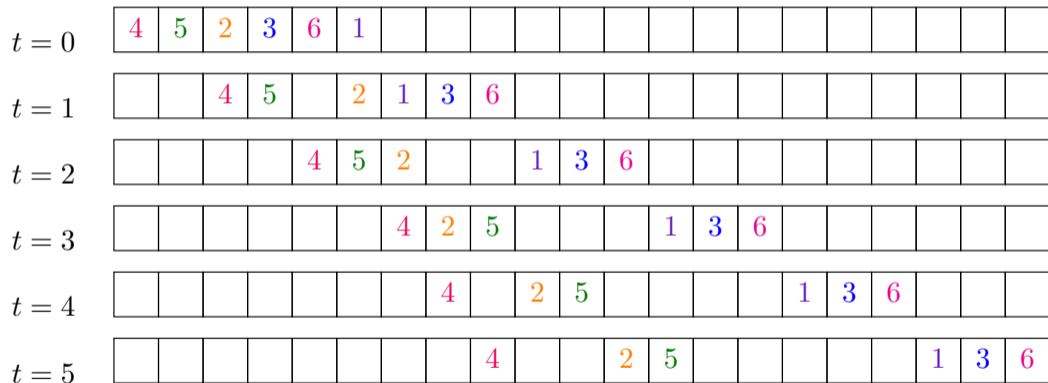
Banff International Research Station for Mathematical Innovation and Discovery (BIRS):  
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## Multicolor box-ball system, Takahashi 1993

A *box-ball system* (BBS) is a dynamical system with balls labeled by numbers 1 through  $n$  in an infinite strip of boxes. Balls take turns jumping to the rightmost empty box, starting with the smallest-numbered ball.



# Box-ball system example ( $t = 0$ through $t = 5$ )



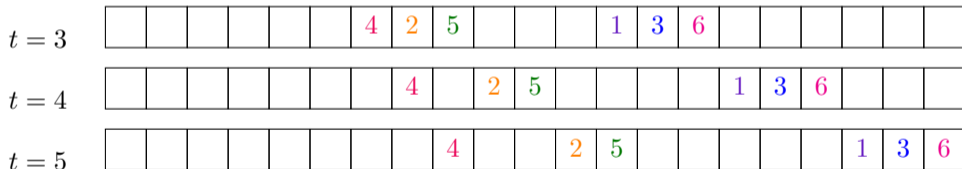
## Solitons and steady state

### Definition

A *soliton* of a box-ball system is an increasing run of balls that moves at a speed equal to their length and is preserved by all future box-ball moves.

### Example

The strings 4, 25, and 136 are solitons:



After a finite number of BBS moves, the system reaches a *steady state* where:

- ▶ the system is decomposed into solitons, i.e., each ball belongs to one soliton
- ▶ the lengths of the solitons are weakly decreasing from right to left

## Tableaux (English notation)

### Definition

A *tableau* is an arrangement of numbers  $\{1, 2, \dots, n\}$  into rows whose lengths are weakly decreasing.

A tableau is *standard* if its rows and columns are increasing.

### Example

Standard Tableaux:

1	2	4
3	5	
6	7	

1	3	6
2	5	
4		

1	3	4
2		
5		
6		

Not a tableau:

1	2	
3	5	4

Nonstandard Tableau:

1	2	3
5	6	7
4		

# Soliton decomposition

## Definition

The *soliton decomposition*  $SD(w)$  of a permutation  $w$  is the tableau whose rows are the solitons stacked from right to left.

## Example

$t = 3$							4	2	5				1	3	6					
$t = 4$							4		2	5				1	3	6				
$t = 5$								4			2	5						1	3	6

$$SD(452361) = \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 2 & 5 & \\ \hline 4 & & \\ \hline \end{array} \text{ with shape } (3, 2, 1).$$

## RSK algorithm

The Robinson–Schensted–Knuth (RSK) insertion algorithm is a bijection

$$w \mapsto (P(w), Q(w))$$

from  $S_n$  onto pairs of size- $n$  standard tableaux of equal shape.

### Example

Let  $w = \mathbf{452361}$ .  $P(w) =$ 

1	3	6
2	5	
4		

 and  $Q(w) =$ 

1	2	5
3	4	
6		

.

(Please ask me if you'd like to see the computation.)

# The $Q$ tableau determines the dynamics of a box-ball system

## Theorem (SUMRY 2021)

If  $Q(\pi) = Q(w)$ , then the box-ball systems of  $\pi$  and  $w$  are identical if we ignore the ball labels, in particular:

- ▶  $\pi$  and  $w$  first reach steady state at the same time, and
- ▶ the soliton decompositions of  $\pi$  and  $w$  have the same shape

## Example

$$\pi = 21435 \text{ and } w = 31425$$

$$Q(\pi) = Q(w) = \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 2 & 4 & \\ \hline \end{array}$$

Both  $\pi$  and  $w$  first reach steady state at  $t = 1$ .

$$SD(\pi) = \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 4 & & \\ \hline 2 & & \\ \hline \end{array} \quad SD(w) = \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 4 & & \\ \hline 3 & & \\ \hline \end{array}$$



## L-shaped soliton decompositions

The time when  $w$  first reaches steady state is called the *time to steady state* of  $w$ .

**Theorem (SUMRY 2021)**

If a permutation has an L-shaped soliton decomposition  $SD = \begin{array}{|c|c|c|} \hline 1 & & \\ \hline & & \\ \hline & & \\ \hline \vdots & & \\ \hline \end{array}, \dots$ , then its time to steady state is either  $t = 0$  or  $t = 1$ .

**Example**

Such permutations include noncrossing involutions and column reading words of standard tableaux.

Both  $\pi = 21435$  and  $w = 31425$  have steady-state time  $t = 1$ .

$$SD(\pi) = \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 4 & & \\ \hline 2 & & \\ \hline \end{array} \quad SD(w) = \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 4 & & \\ \hline 3 & & \\ \hline \end{array}$$

$\pi = 21435 = (12)(34)$  and  $w = 31425$  is the column reading word of  $\begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 3 & 4 & \\ \hline \end{array}$ .

## Maximum steady-state time

### Theorem (UConn 2020)

If  $n \geq 5$  and

$$Q(w) = \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & 4 \\ \hline n & \\ \hline \end{array} \dots \begin{array}{|c|c|} \hline n-2 & n-1 \\ \hline \end{array},$$

then the steady-state time of  $w$  is  $n - 3$ .

### Conjecture

For  $n \geq 4$ , the maximum time to steady state is  $n - 3$ .

### Partial Results (SUMRY 2021):

- ▶ If the shape of  $Q(w)$  is  $(n - 3, 2, 1)$ , the maximum steady-state time is  $n - 3$ .

## When is SD standard?

### Example

$$\text{SD}(452361) = \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 2 & 5 & \\ \hline 4 & & \\ \hline \end{array}$$

$$\text{SD}(21435) = \begin{array}{|c|c|c|} \hline 1 & 3 & 5 \\ \hline 4 & & \\ \hline 2 & & \\ \hline \end{array}$$

$$\text{SD}(31425) = \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 4 & & \\ \hline 3 & & \\ \hline \end{array}$$

### Theorem (UConn 2020)

Given  $w \in S_n$ , the following are equivalent:

1.  $\text{SD}(w)$  is standard
2.  $\text{SD}(w) = P(w)$
3. the shape of  $\text{SD}(w)$  is the same as the shape of  $P(w)$

### Definition

We say that a permutation  $w$  is *good* if the tableau  $\text{SD}(w)$  is standard.

## Good tableaux and Motzkin numbers

**Fact:**  $Q(w)$  determines whether  $w$  is good

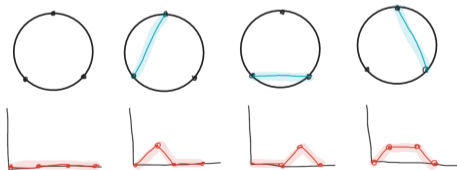
Given a  $Q$ -equivalence class, either all permutations in it are good or all of them are not good.

**Definition (Good tableaux)**

A standard tableau  $T$  is *good* if each permutation whose  $Q$  tableau equals  $T$  is good.

**Conjecture**

$\{Q(w) \mid w \in S_n \text{ and } SD(w) \text{ is standard}\}$  are counted by the Motzkin numbers.



$n = 3$

# Consecutive pattern avoidance

## Definition

A permutation  $\sigma$  is said to be a *consecutive pattern* of another permutation  $w$  if  $w$  has a consecutive subsequence whose elements are in the same relative order as  $\sigma$ .

## Example

$w = 314592687$  contains  $\sigma = 2413$  because the consecutive subsequence 5926 is ordered in the same way as  $\sigma = 2413$ .

## Theorem (SUMRY 2021)

The good permutations are closed under consecutive pattern containment. That is, if a permutation is good, then any consecutive subpermutation is also good.

## Knuth Relations

Suppose  $\pi, w \in S_n$  and  $x < y < z$ .

1.  $\pi$  and  $w$  differ by a Knuth relation of the **first kind** ( $K_1$ ) if

$$\pi = x_1 \dots yxz \dots x_n \text{ and } w = x_1 \dots yzx \dots x_n \text{ or vice versa}$$

2.  $\pi$  and  $w$  differ by a Knuth relation of the **second kind** ( $K_2$ ) if

$$\pi = x_1 \dots xzy \dots x_n \text{ and } w = x_1 \dots zxy \dots x_n \text{ or vice versa}$$

In addition,  $\pi$  and  $w$  differ by a Knuth relation of **both kinds** ( $K_B$ ) if they differ by  $K_1$  and they differ by  $K_2$ , that is,

$$\pi = x_1 \dots y_1 xzy_2 \dots x_n \text{ and } w = x_1 \dots y_1 zxy_2 \dots x_n \text{ or vice versa}$$

where  $x < y_1, y_2 < z$

**Example**

$$326154 \sim^{K_1} 362154$$

$$362154 \sim^{K_B} 362514$$

We say that  $\pi$  and  $w$  are *Knuth equivalent* if they differ by a finite sequence of Knuth relations.

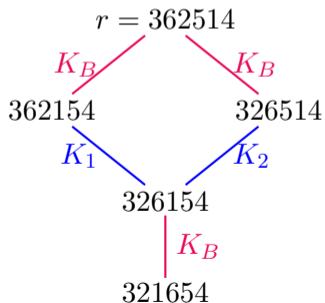
## Facts (Knuth)

- ▶ There is a path of Knuth moves from  $w$  to the row reading word of  $P(w)$ .
- ▶ Two permutations have the same  $P$  tableau if and only if they are in the same Knuth equivalence class.

### Example

The Knuth equivalence class of the row reading word  $r = 362514$  of

1	4
2	5
3	6



## Soliton decompositions and Knuth moves

The soliton decomposition is preserved by non- $K_B$  Knuth moves, but one  $K_B$  move changes the soliton decomposition.

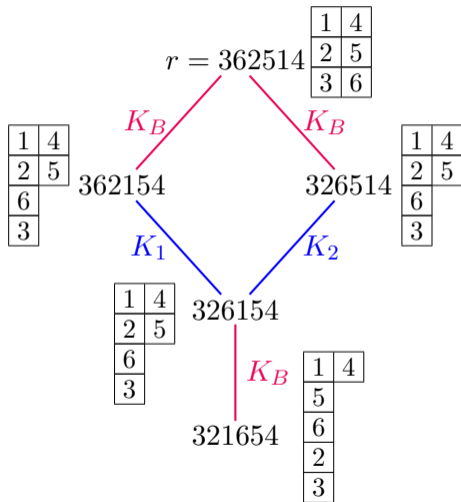
### Theorem (UConn Math REU 2020)

Let  $r$  denote the row reading word of  $P(w)$ .

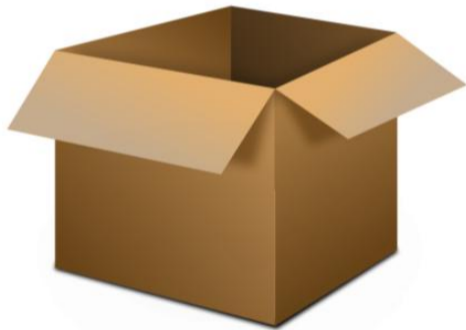
- ▶  $\text{SD}(r) = P(r)$ .
- ▶ If there exists a path of *non- $K_B$*  Knuth moves from  $w$  to  $r$ , then  $\text{SD}(w) = P(w)$ .
- ▶ If there exists a path from  $w$  to  $r$  containing an *odd* number of  $K_B$  moves, then  $\text{SD}(w) \neq P(w)$ .



# Soliton decompositions in the Knuth equivalence class of 362154



Thank you!



## Extra: RSK algorithm example

Let  $w = 452361$ .

$$\begin{array}{r}
 P : \quad 4 \qquad 4 \quad 5 \qquad 2 \quad 5 \qquad 2 \quad 3 \qquad 2 \quad 3 \quad 6 \qquad 1 \quad 3 \quad 6 \\
 \qquad \qquad \qquad \qquad \qquad 4 \qquad \qquad 4 \quad 5 \qquad 4 \quad 5 \qquad \qquad \qquad 2 \quad 5 \qquad \qquad \qquad \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 2 & 5 & \\ \hline 4 & & \\ \hline \end{array} \\
 \\
 Q : \quad 1 \qquad 1 \quad 2 \qquad 1 \quad 2 \qquad 1 \quad 2 \qquad 1 \quad 2 \quad 5 \qquad 1 \quad 2 \quad 5 \\
 \qquad \qquad \qquad \qquad \qquad 3 \qquad \qquad 3 \quad 4 \qquad 3 \quad 4 \qquad \qquad \qquad 3 \quad 4 \qquad \qquad \qquad \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 3 & 4 & \\ \hline 6 & & \\ \hline \end{array}
 \end{array}$$

### Insertion and bumping rule for $P$

Insert  $x$  into the first row of  $P$ . If  $x$  is larger than every element in the first row, add  $x$  to the end of the first row. If not, replace the smallest number larger than  $x$  in row 1 with  $x$ . Insert this number into the row below following the same rules.

### Recording rule for $Q$

For  $Q$ , insert  $1, \dots, n$  in order so that the shape of  $Q$  at each step matches the shape of  $P$ .

## Extra: A localized version of Greene's theorem, part 1

Definition (A localized version of longest  $k$ -increasing subsequences)

Let  $i(u) :=$  the length of a longest increasing subsequence of  $u$ .

For  $w \in S_n$  and  $k \geq 1$ , let  $I_k(w) = \max_{w=u_1|\dots|u_k} \sum_{j=1}^k i(u_j)$ , where the maximum is taken over ways of writing  $w$  as a concatenation  $u_1 | \dots | u_k$  of consecutive subsequences.

### Example

Let  $w = 5623714$ . For short, we write  $I_k := I_k(w)$ . Then

$I_1 = i(w) = 3$  (since the longest increasing subsequences are 567, 237, and 234),

$I_2 = 5$  (witnessed by 56|23714 or 56237|14),

$I_3 = 7$  (witnessed uniquely by 56|237|14), and

$I_k = 7$  for all  $k \geq 3$ .

## Extra: A localized version of Greene's theorem, part 2

Definition (A localized version of longest  $k$ -decreasing subsequences)

Let  $D(u) := 1 + |\{\text{descents of } u\}|$ .

For  $w \in S_n$  and  $k \geq 1$ , let  $D_k(w) = \max_{w=u_1 \sqcup \dots \sqcup u_k} \sum_{j=1}^k D(u_j)$ , where the maximum is taken over ways to write  $w$  as the union of disjoint subsequences  $u_j$  of  $w$ .

### Example

Let  $w = 5623714$ . For short, we write  $D_k := D_k(w)$ . Then

$$D_1 = D(w) = 1 + |\{\text{descents of } 5623714\}| = 1 + |\{2, 5\}| = 3,$$

$D_2 = 6$  (one can take subsequences 531 and 6274, among other partitions),

$D_3 = 7$  (one can take subsequences 52, 631, and 74, among other partitions), and

$D_k = 7$  for all  $k \geq 3$ .

## Extra: A localized version of Greene's theorem, part 3

Theorem (Lewis–Lyu–Pylyavskyy–Sen 2019)

Suppose  $w \in S_n$ . Let  $\Lambda = (\Lambda_1, \Lambda_2, \Lambda_3, \dots)$  denote  $\text{sh SD}(w)$ . Let  $M = (M_1, M_2, M_3, \dots)$  denote the conjugate of  $\Lambda$ . Then, for any  $k$ ,

$$\begin{aligned}I_k(w) &= \Lambda_1 + \Lambda_2 + \dots + \Lambda_k, \\D_k(w) &= M_1 + M_2 + \dots + M_k.\end{aligned}$$

Example

Let  $w = 5623714$ . Then  $\text{sh SD}(w) = (I_1, I_2 - I_1, I_3 - I_2) = (3, 2, 2)$ . We can verify this by computing the soliton decomposition  $\text{SD}(w)$ , which turns out to be the (non-standard) tableau

1	3	4
2	7	
5	6	

.

Note:  $\text{sh SD}(w) = (3, 2, 2)$  is smaller than  $\text{sh } P(w) = (3, 3, 1)$  in the dominance order.

## Extra: Two permutations with L-shaped SD

L-shaped SD which is not a column reading word:

$w = 3217654 = (13)(47)(56)$  is a noncrossing involution.

$$P(w) = Q(w) = \begin{array}{|c|c|} \hline 1 & 4 \\ \hline 2 & 5 \\ \hline 3 & 6 \\ \hline 7 & \\ \hline \end{array} \text{ and } SD(w) = \begin{array}{|c|} \hline 1 & 4 \\ \hline 5 \\ \hline 6 \\ \hline 7 \\ \hline 2 \\ \hline 3 \\ \hline \end{array}$$

An involution which is neither noncrossing nor a column reading word:

$\pi = 5274163 = (15)(37)$  has a crossing.

$$P(\pi) = Q(\pi) = \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 2 & 4 & \\ \hline 5 & 7 & \\ \hline \end{array} \text{ and } SD(\pi) = \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 4 \\ \hline 2 \\ \hline 7 \\ \hline 5 \\ \hline \end{array}$$

## Extra: Good permutations are not closed under classical pattern containment

Starting with  $n = 5$ , a good permutation in  $S_n$  may have a substring which is not good.

### Example

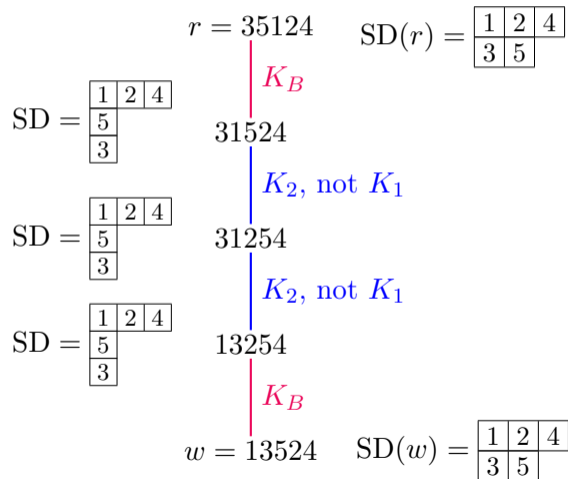
- ▶ The permutation 25143 is good, but its subpermutation 2143 is not good.
- ▶ The permutation 35142 is good, but its subpermutation 3142 is not good.
- ▶ Let  $w = 42513$ , which is a good permutation, and let  $\sigma = 4253$  be a substring of  $w$ . The standardization of  $\sigma$  is 3142, which is not good.

(Therefore, the good permutations cannot be characterized by a set of classical avoided patterns.)



## Extra: Permutations connected by $K_B$ moves and have the same $SD$

Two permutations with the same  $SD$  which are connected by  $K_B$  moves:



## Extra: Solitary waves

(Desmos link)

