Box-Ball Systems and Robinson–Schensted–Knuth Tableaux

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Slides available at egunawan.github.io/talks/21/december/lancaster21/main.pdf

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Solitary waves (solitons)

Scott Russell's first encounter of solitary waves at the Union Canal (near Edinburgh):

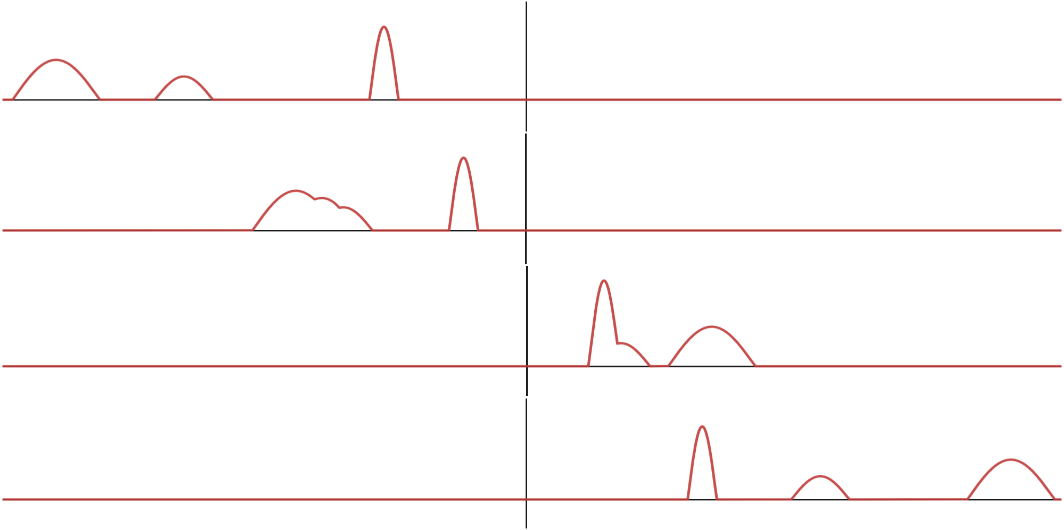
'I was observing the motion of a boat which was rapidly drawn along a narrow channel by a pair of horses, when the boat suddenly stopped—not so the mass of water in the channel which it had put in motion; it accumulated round the prow of the vessel in a state of violent agitation, then suddenly leaving it behind, rolled forward with great velocity, assuming the form of a large solitary elevation, a rounded, smooth and well-defined heap of water, which continued its course along the channel apparently **without change of form or diminution of speed**. I followed it on horseback, and overtook it still rolling on at a rate of some eight or nine miles an hour, preserving its original figure some thirty feet long and a foot to a foot and a half in height. Its height gradually diminished, and after a chase of one or two miles I lost it in the windings of the channel. Such, in the month of August 1834, was my first chance interview with that singular and beautiful phenomenon which I have called the Wave of Translation.'

Soliton on the Scott Russell Aqueduct on the Union Canal near Heriot-Watt University, July 1995

Credit: ma.hw.ac.uk/solitons/press.html

Solitary waves

(Desmos link)



Multicolor box-ball system (BBS), Takahashi 1993

A *box-ball system* (BBS) is a dynamical system of BBS configurations.

- At each configuration, balls are labeled by numbers 1 through n in an infinite strip of boxes.
- Each box can fit at most one ball.

Example

A possible BBS configuration:

$\cdots$	4	5	2	3	6	1									$\cdots$

Box-ball move (from $t = 0$ to $t = 1$)

Balls take turns jumping to the first (leftmost) empty box, starting with the smallest-numbered ball.

$t = 0$	4	5	2	3	6	1			
	4	5	2	3	6		1		
	4	5		3	6	2	1		
	4	5			6	2	1	3	
		5	4		6	2	1	3	
			4	5	6	2	1	3	
$t = 1$			4	5		2	1	3	6

Box-ball moves ($t = 0$ through $t = 5$)

$t = 0$	4	5	2	3	6	1															
$t = 1$			4	5		2	1	3	6												
$t = 2$					4	5	2			1	3	6									
$t = 3$							4	2	5				1	3	6						
$t = 4$								4		2	5					1	3	6			
$t = 5$									4			2	5						1	3	6

Solitons and steady state

Definition

A *soliton* of a box-ball system is an increasing run of balls that moves at a speed equal to its length and is preserved by all future box-ball moves.

Example

The strings 4, 25, and 136 are solitons:

$t = 3$							4	2	5				1	3	6						
$t = 4$								4		2	5					1	3	6			
$t = 5$									4			2	5						1	3	6

After a finite number of BBS moves, the system reaches a *steady state* where:

- the system is decomposed into solitons, i.e., each ball belongs to one soliton
- the lengths of the solitons are weakly decreasing from right to left

Tableaux (English notation)

Definition

- A *tableau* is an arrangement of numbers $\{1, 2, \ldots, n\}$ into rows whose lengths are weakly decreasing.
- A tableau is *standard* if its rows and columns are increasing.

Example

Standard Tableaux:

1	2	4
3	5	
6	7	

1	3	6
2	5	
4		

1	3	4
2		
5		
6		

Not a tableau:

1	2	
3	5	4

Nonstandard Tableau:

1	2	3
5	6	7
4		

Soliton decomposition

Definition

To construct *soliton decomposition* $\mathrm{SD}(w)$ of a permutation w, stack solitons so that the rightmost soliton is placed on the first row, the soliton to its left is placed on the second row, and so on.

Example

$t = 3$							4	2	5				1	3	6						
$t = 4$								4		2	5					1	3	6			
$t = 5$									4			2	5						1	3	6

$\mathrm{SD}(452361) =$

1	3	6
2	5	
4		

with shape $(3, 2, 1)$.

RSK algorithm

The Robinson–Schensted–Knuth (RSK) insertion algorithm is a bijection

$$\pi \mapsto (\mathrm{P}(\pi), \mathrm{Q}(\pi))$$

from S_n onto pairs of size-n standard tableaux of equal shape.

Example

Let $w =$ **452361**. $\mathrm{P}(w) = \begin{array}{|c|c|c|}\hline 1 & 3 & 6 \\ \hline 2 & 5 \\ \cline{1-2} 4 \\ \cline{1-1}\end{array}$ and $\mathrm{Q}(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 3 & 4 \\ \cline{1-2} 6 \\ \cline{1-1}\end{array}$.

RSK algorithm example

Let $w = \mathbf{452361}$.

$$\text{P}: \quad \mathbf{4} \qquad \begin{array}{cc} 4 & \mathbf{5} \end{array} \qquad \begin{array}{cc} \mathbf{2} & 5 \\ 4 & \end{array} \qquad \begin{array}{cc} 2 & \mathbf{3} \\ 4 & 5 \end{array} \qquad \begin{array}{ccc} 2 & 3 & \mathbf{6} \\ 4 & 5 & \end{array} \qquad \begin{array}{ccc} \mathbf{1} & 3 & 6 \\ 2 & 5 & \\ 4 & & \end{array} \qquad \text{P}(w) = \begin{array}{|c|c|c|} \hline 1 & 3 & 6 \\ \hline 2 & 5 & \\ \cline{1-2} 4 & & \\ \cline{1-1} \end{array}$$

$$\text{Q}: \quad \mathbf{1} \qquad \begin{array}{cc} 1 & \mathbf{2} \end{array} \qquad \begin{array}{cc} 1 & 2 \\ \mathbf{3} & \end{array} \qquad \begin{array}{cc} 1 & 2 \\ 3 & \mathbf{4} \end{array} \qquad \begin{array}{ccc} 1 & 2 & \mathbf{5} \\ 3 & 4 & \end{array} \qquad \begin{array}{ccc} 1 & 2 & 5 \\ 3 & 4 & \\ \mathbf{6} & & \end{array} \qquad \text{Q}(w) = \begin{array}{|c|c|c|} \hline 1 & 2 & 5 \\ \hline 3 & 4 & \\ \cline{1-2} 6 & & \\ \cline{1-1} \end{array}$$

Insertion and bumping rule for P

- Insert x into the first row of P.
- If x is larger than every element in the first row, add x to the end of the first row.
- If not, replace the smallest number larger than x in row 1 with x. Insert this number into the row below following the same rules.

Recording rule for Q

For Q, insert $1, \ldots, n$ in order so that the shape of Q at each step matches the shape of P.

The Q tableau determines the dynamics of a box-ball system

Theorem (SUMRY 2021)

If $\mathrm{Q}(\pi) = \mathrm{Q}(w)$, then the box-ball systems of π and w are identical if we ignore the ball labels, in particular:

- π and w first reach steady state at the same time, and
- the soliton decompositions of π and w have the same shape

Example

$$\pi = 21435 \text{ and } w = 31425$$

$$\mathrm{Q}(\pi) = \mathrm{Q}(w) = \begin{array}{|c|c|c|}\hline 1 & 3 & 5 \\ \hline 2 & 4 \\ \cline{1-2} \end{array}$$

Both π and w first reach steady state at $t = 1$.

$$\mathrm{SD}(\pi) = \begin{array}{|c|c|c|}\hline 1 & 3 & 5 \\ \hline 4 \\ \cline{1-1} 2 \\ \cline{1-1} \end{array} \qquad \mathrm{SD}(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 4 \\ \cline{1-1} 3 \\ \cline{1-1} \end{array}$$

Questions (steady-state time)

The time when a permutation w first reaches steady state is called the *steady-state time* of w.

- Given a Q tableau, find a formula to compute the steady-state time for all permutations in the Q-tableau class.
- Find an upper bound for steady-state time.

L-shaped soliton decompositions

Theorem (SUMRY 2021)

If a permutation has an L-shaped soliton decomposition $\mathrm{SD} = \begin{array}{|c|c|c|l}\hline 1 & & & \cdots \\ \hline & \multicolumn{3}{l}{} \\ \cline{1-1} & \multicolumn{3}{l}{} \\ \cline{1-1} \multicolumn{1}{c}{\vdots} & \multicolumn{3}{l}{} \end{array}$,

then its steady-state time is either $t = 0$ or $t = 1$.

Example

Such permutations include noncrossing involutions and column reading words of standard tableaux.

Both $\pi = 21435$ and $w = 31425$ have steady-state time $t = 1$.

$$\mathrm{SD}(\pi) = \begin{array}{|c|c|c|}\hline 1 & 3 & 5 \\ \hline 4 & \multicolumn{2}{c}{} \\ \cline{1-1} 2 & \multicolumn{2}{c}{} \\ \cline{1-1} \end{array} \qquad \mathrm{SD}(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 4 & \multicolumn{2}{c}{} \\ \cline{1-1} 3 & \multicolumn{2}{c}{} \\ \cline{1-1} \end{array}$$

$\pi = 21435 = (12)(34)$ and $w = 31425$ is the column reading word of $\begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 3 & 4 & \multicolumn{1}{c}{} \\ \cline{1-2} \end{array}$.

Maximum steady-state time

Theorem (UConn 2020)

If $n \geq 5$ and

$$Q(w) = \begin{array}{|c|c|}\hline 1 & 2 \\ \hline 3 & 4 \\ \hline n & \\ \cline{1-1} \end{array} \cdots \begin{array}{|c|c|}\hline n-2 & n-1 \\ \hline \end{array},$$

then the steady-state time of w is $n-3$.

Conjecture

For $n \geq 4$, the steady-state time of a permutation in S_n is at most $n-3$.

Partial Result (SUMRY 2021)

If the shape of $Q(w)$ is $(n-3, 2, 1)$, the steady-state time is at most $n-3$.

Box-Ball System Example ($t = 0$ through 5)

Let $w = 452361$. Then $Q(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 3 & 4 \\ \cline{1-2} 6 \\ \cline{1-1}\end{array}$ and the steady-state time of w is $3 = n - 3$.

$t = 0$	4	5	2	3	6	1															
$t = 1$			4	5		2	1	3	6												
$t = 2$					4	5	2			1	3	6									
$t = 3$							4	2	5				1	3	6						
$t = 4$								4		2	5					1	3	6			
$t = 5$									4			2	5						1	3	6

Questions (soliton decomposition)

- When is the soliton decomposition SD a standard tableau?
- Can we classify permutations with standard SD using pattern avoidance?

When is SD standard?

Example

$$\mathrm{SD}(452361) = \begin{array}{|c|c|c|}\hline 1 & 3 & 6 \\ \hline 2 & 5 \\ \cline{1-2} 4 \\ \cline{1-1}\end{array} \qquad \mathrm{SD}(21435) = \begin{array}{|c|c|c|}\hline 1 & 3 & 5 \\ \hline 4 \\ \cline{1-1} 2 \\ \cline{1-1}\end{array} \qquad \mathrm{SD}(31425) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 4 \\ \cline{1-1} 3 \\ \cline{1-1}\end{array}$$

Theorem (UConn 2020)

Given $w \in S_n$, the following are equivalent:

1. $\mathrm{SD}(w)$ is standard
2. $\mathrm{SD}(w) = \mathrm{P}(w)$
3. the shape of $\mathrm{SD}(w)$ is equal to the shape of $\mathrm{P}(w)$

Definition

We say that a permutation w is *good* if the tableau $\mathrm{SD}(w)$ is standard.

$\mathrm{Q}(w)$ determines whether w is good

Proposition

Given a Q-equivalence class, either all permutations in it are good or all of them are not good.

Proof

1. The tableau $\mathrm{Q}(w)$ determines the shape of $\mathrm{SD}(w)$.
2. $\mathrm{SD}(w)$ is standard iff $\operatorname{sh}\mathrm{SD}(w) = \operatorname{sh}\mathrm{P}(w)$

Suppose $\mathrm{Q}(w) = \mathrm{Q}(\pi)$. Then

$$\begin{aligned}\mathrm{SD}(w) \text{ is standard} &\Longrightarrow \operatorname{sh}\mathrm{SD}(\pi) = \operatorname{sh}\mathrm{SD}(w) = \operatorname{sh}\mathrm{P}(w) = \operatorname{sh}\mathrm{P}(\pi) \\ &\Longrightarrow \mathrm{SD}(\pi) \text{ is standard, that is, } \pi \text{ is also good}\end{aligned}$$

Definition (Good tableaux)

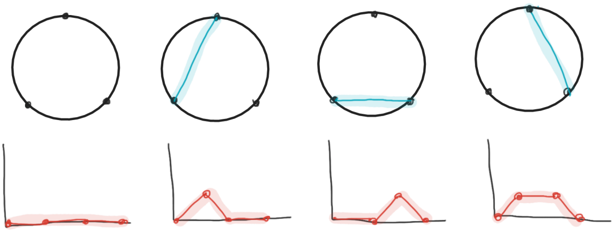
A standard tableau T is *good* if each permutation whose Q tableau equals T is good.

Good tableaux and Motzkin numbers

Conjecture

The good standard tableaux, $\{\mathrm{Q}(w) \mid w \in S_n \text{ and } \mathrm{SD}(w) \text{ is standard}\}$, are counted by the Motzkin numbers.

Some objects counted by Motzkin numbers:



$n = 3$

The first few Motzkin numbers are 1, 2, 4, 9, 21, 51, 127, 323, 835.

Consecutive pattern avoidance

Lemma (SUMRY 2021)

If a standard tableau T is good, then the tableau T' obtained by removing the largest k cells from T is also good.

Example

$$T = \begin{array}{|c|c|c|c|}\hline 1 & 2 & 3 & 7 \\ \hline 4 & 5 & 6 \\ \cline{1-3} 8 \\ \cline{1-1}\end{array} \text{ is good.}$$

$$T' = \begin{array}{|c|c|c|}\hline 1 & 2 & 3 \\ \hline 4 & 5 \\ \cline{1-2}\end{array} \text{ is also good.}$$

Consecutive pattern avoidance

Definition

A permutation σ is a *consecutive pattern* of another permutation w if w has a consecutive subsequence whose elements are in the same relative order as σ.

Example

$w = 314592687$ contains $\sigma = 2413$ because the consecutive subsequence 5926 is ordered in the same way as $\sigma = 2413$.

Theorem (SUMRY 2021)

The good permutations are closed under consecutive pattern containment. That is, if a permutation is good, then any consecutive subpermutation is also good.

Corollary

The good permutations can be characterized by a (possibly infinite) set of consecutive avoided patterns.

Knuth Relations

Suppose π, $w \in S_n$ and $x < y < z$.

1. π and w differ by a Knuth relation of the **first kind** (K_1) if
$$\pi = x_1 \dots yxz \dots x_n \text{ and } w = x_1 \dots yzx \dots x_n \text{ or vice versa}$$
2. π and w differ by a Knuth relation of the **second kind** (K_2) if
$$\pi = x_1 \dots xzy \dots x_n \text{ and } w = x_1 \dots zxy \dots x_n \text{ or vice versa}$$

In addition, π and w differ by a Knuth relation of **both kinds** (K_B) if they differ by K_1 and they differ by K_2, that is,

$$\pi = x_1 \dots y_1 xzy_2 \dots x_n \text{ and } w = x_1 \dots y_1 zxy_2 \dots x_n \text{ or vice versa}$$

where $x < y_1, y_2 < z$

Example $\quad \mathbf{326}154 \sim^{K_1} \mathbf{362}154 \qquad 362\mathbf{154} \sim^{K_B} 362\mathbf{514}$

We say that π and w are *Knuth equivalent* if they differ by a finite sequence of Knuth relations.

P-tableaux and Knuth moves

Theorem (Knuth, 1970)

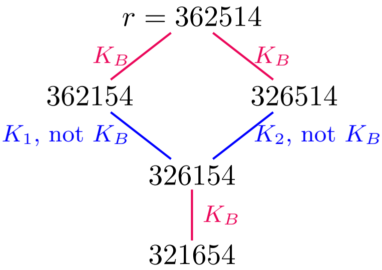
- *There is a path of Knuth moves from w to the row reading word of $P(w)$.*
- *Two permutations have the same P tableau iff they are in the same Knuth equivalence class.*

Example

The Knuth equivalence class of the row reading word $r = 362514$ of

1	4
2	5
3	6

:



Soliton decompositions and Knuth moves

The soliton decomposition is preserved by non-K_B Knuth moves, but one K_B move changes the soliton decomposition.

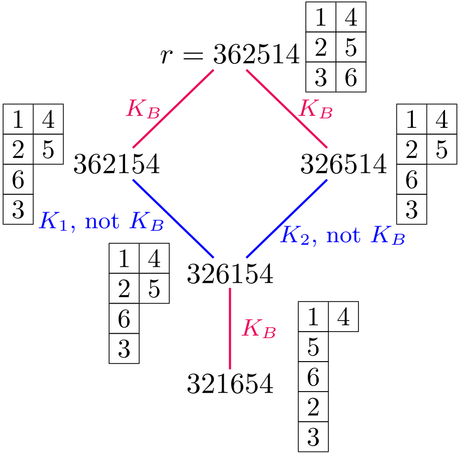
Theorem (UConn Math REU 2020)

Let r denote the row reading word of $\mathrm{P}(w)$.

- If there exists a path of *non-*K_B Knuth moves from w to r, then $\mathrm{SD}(w) = \mathrm{P}(w)$. In particular, $\mathrm{SD}(r) = \mathrm{P}(r)$.
- If there exists a path from w to r containing an *odd* number of K_B moves, then $\mathrm{SD}(w) \neq \mathrm{P}(w)$.

Example

Soliton decompositions of the Knuth equivalence class of 362154:



Thank you!

Y	O	U	!
A	N	K	
T	H		

Greene's theorem, slide 1/3

Definition (longest k-increasing subsequences)

A subsequence σ of w is called *k-increasing* if, as a set, it can be written as a disjoint union

$$\sigma = \sigma_1 \sqcup \sigma_2 \sqcup \cdots \sqcup \sigma_k$$

where each σ_i is an increasing subsequence of w. Let $\mathrm{i}_k := \mathrm{i}_k(w)$ denote the length of a longest k-increasing subsequence of w.

Example (Let $w = 5623714$.)

- The longest 1-increasing subsequences are 567, 237, and 234.
- The longest 2-increasing subsequence is given by $562374 = 567 \sqcup 234$.
- A longest 3-increasing subsequence (among others) is given by $5623714 = 56 \sqcup 237 \sqcup 14$.
- Thus,

$$\mathrm{i}_1 = 3, \qquad \mathrm{i}_2 = 6, \qquad \text{and} \qquad \mathrm{i}_k = 7 \text{ if } k \geq 3.$$

Greene's theorem, slide 2/3

Definition (longest k-decreasing subsequences)

Similarly, a subsequence σ of w is called *k-decreasing* if, as a set, it can be written as a disjoint union

$$\sigma = \sigma_1 \sqcup \sigma_2 \sqcup \cdots \sqcup \sigma_k$$

where each σ_i is an decreasing subsequence of w. Let $\mathrm{d}_k := \mathrm{d}_k(w)$ denote the length of a longest k-decreasing subsequence of w.

Example (Let $w = 5623714$.)

- The longest 1-decreasing subsequences are 521, 621, 531, and 631.
- A longest 2-decreasing subsequence (among others) is given by $52714 = 521 \sqcup 74$.
- A longest 3-decreasing subsequence (among others) is given by $5623714 = 52 \sqcup 631 \sqcup 74$.
- Thus,

$$\mathrm{d}_1 = 3, \qquad \mathrm{d}_2 = 5, \qquad \text{and} \qquad \mathrm{d}_k = 7 \ \text{ if } \ k \geq 3.$$

Greene's theorem, slide 3/3

Theorem (Greene, 1974)

Suppose $w \in S_n$*. Let* $\lambda = (\lambda_1, \lambda_2, \lambda_3, \dots)$ *denote the RS partition of* w*, that is, let* $\lambda = \operatorname{sh} P(w)$*. Let* $\mu = (\mu_1, \mu_2, \mu_3, \dots)$ *denote the conjugate of* λ*. Then, for any* k*,*

$$\mathrm{i}_k(w) = \lambda_1 + \lambda_2 + \ldots + \lambda_k,$$
$$\mathrm{d}_k(w) = \mu_1 + \mu_2 + \ldots + \mu_k.$$

Example

By Greene's theorem, the RS partition is equal to $\lambda = (\mathrm{i}_1, \mathrm{i}_2 - \mathrm{i}_1, \mathrm{i}_3 - \mathrm{i}_2) = (3, 3, 1)$. We can verify this by computing the RS tableaux

$$P(w) = \begin{array}{|c|c|c|}\hline 1 & 3 & 4 \\ \hline 2 & 6 & 7 \\ \hline 5 \\ \cline{1-1}\end{array}, \qquad Q(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 5 \\ \hline 3 & 4 & 7 \\ \hline 6 \\ \cline{1-1}\end{array}.$$

A localized version of Greene's theorem, slide 1/3

Definition (A localized version of longest k-increasing subsequences)

Let $\mathrm{i}(u) :=$ the length of a longest increasing subsequence of u.

For $w \in S_n$ and $k \geq 1$, let $\mathrm{I}_k(w) = \max_{w = u_1 | \cdots | u_k} \sum_{j=1}^{k} \mathrm{i}(u_j)$, where the maximum is taken over ways of writing w as a concatenation $u_1 \mid \cdots \mid u_k$ of consecutive subsequences.

Example

Let $w = 5623714$. For short, we write $\mathrm{I}_k := \mathrm{I}_k(w)$. Then

$\mathrm{I}_1 = \mathrm{i}(w) = 3$ (since the longest increasing subsequences are 567, 237, and 234),

$\mathrm{I}_2 = 5$ (witnessed by $56|23714$ or $56237|14$),

$\mathrm{I}_3 = 7$ (witnessed uniquely by $56|237|14$), and

$\mathrm{I}_k = 7$ for all $k \geq 3$.

A localized version of Greene's theorem, slide 2/3

Definition (A localized version of longest k-decreasing subsequences)

Let $\mathrm{D}(u) \coloneqq 1 + |\{\text{descents of } u\}|$.

For $w \in S_n$ and $k \geq 1$, let $\mathrm{D}_k(w) = \max_{w = u_1 \sqcup \cdots \sqcup u_k} \sum_{j=1}^{k} \mathrm{D}(u_j)$, where the maximum is taken over ways to write w as the union of disjoint subsequences u_j of w.

Example

Let $w = 5623714$. For short, we write $\mathrm{D}_k \coloneqq \mathrm{D}_k(w)$. Then

$\mathrm{D}_1 = \mathrm{D}(w) = 1 + |\text{descents of } 5623714| = 1 + |\{2, 5\}| = 3$,

$\mathrm{D}_2 = 6$ (one can take subsequences 531 and 6274, among other partitions),

$\mathrm{D}_3 = 7$ (one can take subsequences 52, 631, and 74, among other partitions), and

$\mathrm{D}_k = 7$ for all $k \geq 3$.

A localized version of Greene's theorem, slide 3/3

Theorem (Lewis–Lyu–Pylyavskyy–Sen 2019)

Suppose $w \in S_n$. *Let* $\Lambda = (\Lambda_1, \Lambda_2, \Lambda_3, \dots)$ *denote* $\operatorname{sh} \mathrm{SD}(w)$. *Let* $M = (M_1, M_2, M_3, \dots)$ *denote the conjugate of* Λ. *Then, for any* k,

$$\mathrm{I}_k(w) = \Lambda_1 + \Lambda_2 + \dots + \Lambda_k,$$
$$\mathrm{D}_k(w) = M_1 + M_2 + \dots + M_k.$$

Example

Let $w = 5623714$. By the above theorem, $\operatorname{sh} \mathrm{SD}(w) = (\mathrm{I}_1, \mathrm{I}_2 - \mathrm{I}_1, \mathrm{I}_3 - \mathrm{I}_2) = (3, 2, 2)$. We can verify this by computing the soliton decomposition $\mathrm{SD}(w)$, which turns out to be the (non-standard) tableau

1	3	4
2	7	
5	6	

.

Note: $\operatorname{sh} \mathrm{SD}(w) = (3, 2, 2)$ is smaller than $\operatorname{sh} P(w) = (3, 3, 1)$ in the dominance order.

Examples: permutations with L-shaped SD

A permutation with L-shaped SD which is not a column reading word:

$w = 3217654 = (13)(47)(56)$ is a noncrossing involution.

$$\mathrm{P}(w) = \mathrm{Q}(w) = \begin{array}{|c|c|}\hline 1 & 4 \\ \hline 2 & 5 \\ \hline 3 & 6 \\ \hline 7 \\ \cline{1-1}\end{array} \qquad \text{and} \qquad \mathrm{SD}(w) = \begin{array}{|c|c|}\hline 1 & 4 \\ \hline 5 \\ \cline{1-1} 6 \\ \cline{1-1} 7 \\ \cline{1-1} 2 \\ \cline{1-1} 3 \\ \cline{1-1}\end{array}$$

An involution which is neither noncrossing nor a column reading word:

$\pi = 5274163 = (15)(37)$ has a crossing.

$$\mathrm{P}(\pi) = \mathrm{Q}(\pi) = \begin{array}{|c|c|c|}\hline 1 & 3 & 6 \\ \hline 2 & 4 \\ \cline{1-2} 5 & 7 \\ \cline{1-2}\end{array} \qquad \text{and} \qquad \mathrm{SD}(\pi) = \begin{array}{|c|c|c|}\hline 1 & 3 & 6 \\ \hline 4 \\ \cline{1-1} 2 \\ \cline{1-1} 7 \\ \cline{1-1} 5 \\ \cline{1-1}\end{array}$$

Good permutations are not closed under classical pattern containment

Starting with $n = 5$, a good permutation in S_n may have a substring which is not good.

Example

- The permutation 25143 is good, but its subpermutation 2143 is not good.
- The permutation 35142 is good, but its subpermutation 3142 is not good.
- Let $w = 42513$, which is a good permutation, and let $\sigma = 4253$ be a substring of w. The standardization of σ is 3142, which is not good.

(Therefore, the good permutations cannot be characterized by a set of classical avoided patterns.)

Permutations connected by K_B moves and have the same SD

Two permutations with the same SD which are connected by K_B moves:

$r = 35124$ $\quad \mathrm{SD}(r) = \begin{array}{|c|c|c|}\hline 1 & 2 & 4 \\ \hline 3 & 5 & \\ \cline{1-2} \end{array}$

K_B

31524 $\quad \mathrm{SD} = \begin{array}{|c|c|c|}\hline 1 & 2 & 4 \\ \hline 5 \\ \cline{1-1} 3 \\ \cline{1-1} \end{array}$

K_2, not K_1

31254 $\quad \mathrm{SD} = \begin{array}{|c|c|c|}\hline 1 & 2 & 4 \\ \hline 5 \\ \cline{1-1} 3 \\ \cline{1-1} \end{array}$

K_2, not K_1

13254 $\quad \mathrm{SD} = \begin{array}{|c|c|c|}\hline 1 & 2 & 4 \\ \hline 5 \\ \cline{1-1} 3 \\ \cline{1-1} \end{array}$

K_B

$w = 13524$ $\quad \mathrm{SD}(w) = \begin{array}{|c|c|c|}\hline 1 & 2 & 4 \\ \hline 3 & 5 & \\ \cline{1-2} \end{array}$