Sec 3.6 Normalizers (and: another application of cosets!) not normal in G I dea: . If H < G but H A G we want to measure how far H is from being normal. · Recall the def: HJG (def) xH=Hx V xEG "One way to measure how far It is from being normal is to check how may elt XEG satisfy XH=HX - Think of each elt x & G as voting yes or no to the normality of H vote of x Yes if ×H=H× if ×H ≠H× · Remitvery XEH votes YES. Why? XEH => XH=H (Sec 3.3) · Every XEG votes YES iff HJG . If H is not normal, there is at least one elt voting no. Def (for the etts XEG which vote Yes in favor of H's normality) The normalizer of H in G, denoted NG (H), is the set { x E G | x H = Hx } $= \{x \in G \mid x H x' = H\}$ Slides 3.3 "Normal subgroups" Wording We say NG(H) is the set of elements that normalize H. Prop 1 If $x \in N_{\mathcal{C}}(\mathcal{H})$, then $x \mathcal{H} \subseteq N_{\mathcal{G}}(\mathcal{H})$. Lemma (Prop 2 from Slides 3.2 "Cosets") • × +1 = y H for all y = x H (So it does not matter which coset representative) you choose, × or y · Hx = Hy for all y ∈ Hx for the same reason Pf of Suppose x & NG(H). Then xH=Hx by def of NG(H). Prop 1 Let y E x H. (Think to self: my goal is to show y E NG (H), ie I want to show y H = Hy Then yH = xH by above lemma = Hx by (*) = Hy by above lemma Rem Prop 1 tells us that members of a left coset vote together as a block: members of xH <u>all</u> vote yes (when xH = Hx) or <u>all</u> vote no (when XH ≠ HX)

$$\frac{\operatorname{Prop} 2}{|K_{G}(H)|} \text{ is a multiple of } [H].$$

$$\frac{\operatorname{Pr}}{|H|} \operatorname{Prop} 1 \text{ tells us that } N_{G}(H) \text{ consists of entire left cosets of H (at [ast one, H itself). From Sobles 3.2 'Goeds'' use know the left cosets are the same size and algorith. Carton example: Partitions of G by the right cosets by the left cosets by the right cosets
$$\frac{\operatorname{Pr}}{|H|} \underbrace{H_{H}}_{HW} \xrightarrow{H_{H}}_{HW} \xrightarrow{H_{H}}_{HW} \xrightarrow{H_{H}}_{HW}$$

$$The elts of the cosets H and xtlettx all vote Yes. The elts of the left coset gH all vote No since $gH \neq Hg$.
$$\frac{tH}{H} \xrightarrow{H_{H}}_{HW} \xrightarrow{H_{H}}_{HW}$$

$$The entrailizer of H in G is N_{G}(H) = H \cup xH$$

$$\operatorname{Rem} \quad \text{The two extreme cases for } N_{G}(H) = \operatorname{are}:$$

$$N_{G}(H) = G \quad \text{iff } H \triangleleft G$$

$$N_{G}(H) = H \quad (when H is as far from normal as possible) \xrightarrow{H_{H}}_{H} \xrightarrow{H_{H}}_{H} = \operatorname{and}_{H} \xrightarrow{H_{H}}_{H} \xrightarrow{H_{H$$$$$$

So
$$N_{D_{c}}(\langle f \rangle) = \langle f \rangle \cup r^{3} \langle f \rangle$$

 $= \{e, f, r^{3}, r^{3} f \rangle$
 $= V_{4}$
 $r^{3} = r^{3} f$

Exercise

Find a pattern for
$$N_{Dn}(\langle f \rangle)$$
 "it depends on whether "
n is even (odd
Thm
Let $H < G$. Then $N_G(H) < G$.
Pf (we need to check all 3 properties of being a subgroup)
Recall $N_G(H) = f \times e G | \times H \times^{-1} = H f$.
- Contains e $eHe^{-1} = f ehe^{-1} | h \in H f$
= H
- Inverses exist Let $x \in N_G(H)$. We need to show $x^{-1} \in N_G(H)$
Then $x H \times^{-1} \stackrel{\text{(#)}}{=} H$
So $x^{-1} H \langle e^{-1} | e^{-1} = x^{-1} H \times (replace (x^{-1})^{-1} with x)$
 $= x^{-1} (x H \times^{-1}) \times (by (H))$
 $= eHe$
 $= H$

$$\frac{\text{operation of } G}{\text{Suppose } x, y \in N_G(H), \text{ meaning}} \times H_{x^{-1}} \stackrel{(\#)}{=} H \text{ and } gH_{y^{-1}} \stackrel{(\#)}{=} H.$$

$$We \text{ need to show } x y \in N_G(H), \text{ meaning } x gH(xy)^{-1} = H.$$

$$But x gH(xy)^{-1} = x gH_y^{-1} x^{-1} \qquad \begin{pmatrix} \text{"shoes-socks property" of } \\ \text{Thverses: } (xy)^{-1} = y^{-1} x^{-1} \end{pmatrix}$$

$$= x H x^{-1} \qquad (by (\#))$$

$$= H_{\Box} \qquad (by (\# x))$$

$$Thm$$

Every subgroup of G is normal in its normalizer in G: If H < G, then $H < N_G(H) < G$ <u>Pf</u> Let $H < N_G(H)$. To show $H < N_G(H)$, we need to show $xH=Hx \forall x\in N_G(H)$.

But, ×H=H× ¥ × ∈ NG(#) by def.